

LPC  
3pm-8pm  
6-7.30

# 37181 DISCRETE MATHEMATICS

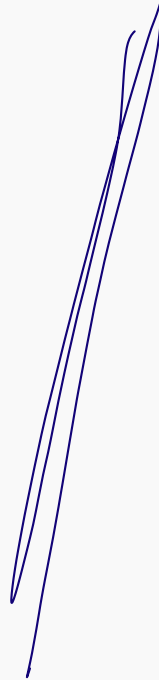
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Lecture 1: propositional logic

# PLAN

- introduction, subject outline
- truth tables
- logical equivalence
- tautology
- logic circuits



# INTRODUCTION

- please read the Subject Outline

- two lectures each week:

- one live zoom lecture Monday 1pm-2:20pm;
- second recorded lecture to watch anytime Tuesday.

- tutorial - online for overseas-based students; on-campus for Sydney-based students. Both running as “whiteboard workshops”. Marks for “active participation” each week. Starts **this week**

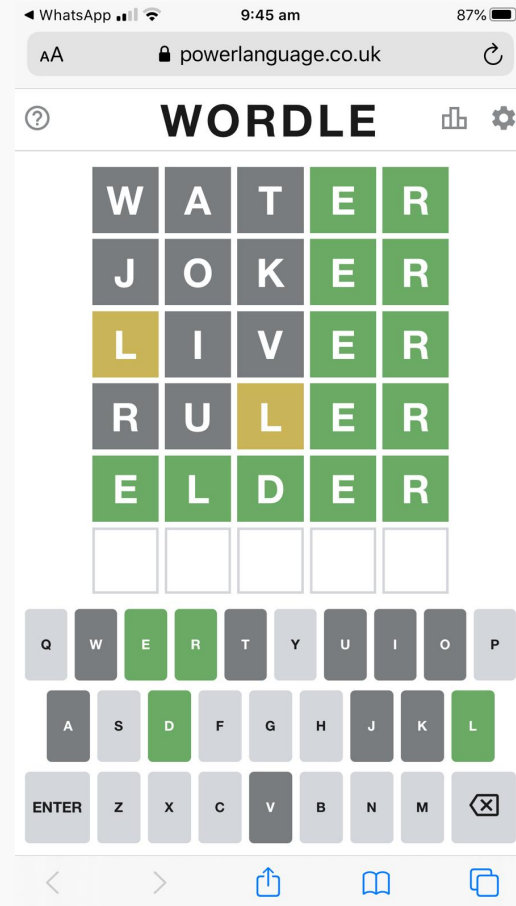
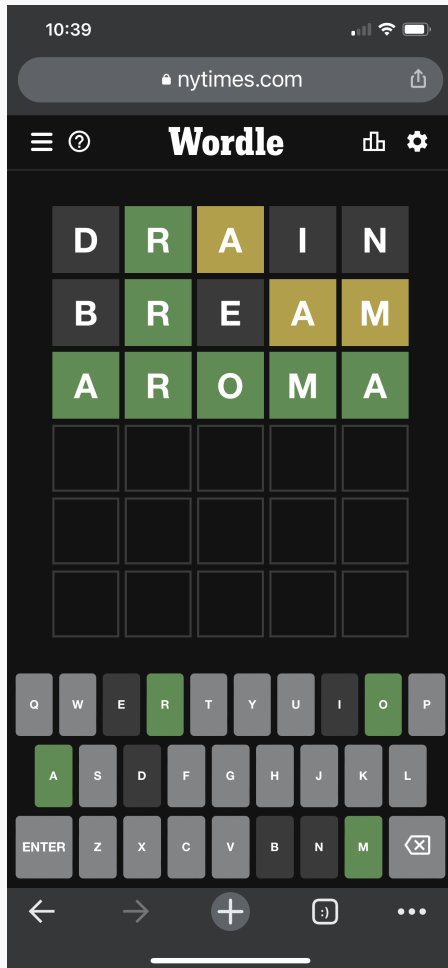
- Learning Progress Checks (quizzes) (LPCs) online Tuesdays 3pm-8pm starting Week 2. Dates in Subject Outline

# WHITEBOARD WORKSHOPS





# CHEATING



## Definition

A statement is a sentence that can (theoretically) be assigned a value of *true* or *false*.

# LOGIC

## Definition

A statement is a sentence that can (theoretically) be assigned a value of *true* or *false*.

Eg:

1. Um, like, whatever NO
2. All positive integers are prime YES (False!)
3. In 2019, all live lectures were recorded at UTS YES False
4. In the year 4000BC, at the location of UTS Building 1, it was raining on the 5th of March at 10am YES
5. When will this lecture end? NO.



6.

YES

# LOGICAL CONNECTIVES

We can build up more complicated statements out of simpler ones using *logical connectives* like *and* and *or*.

Eg:

1. Murray has neat handwriting and Murray has long hair.
2. Murray has neat handwriting or Murray has long hair.
3. Murray does not have neat handwriting and Murray has short hair.

not (Murray has long hair)

## PRECISE MEANING: TRUTH TABLE

English (or any natural human language) can be imprecise, so instead of using our “*intuition*” we define what “and” and “or” and “not” mean using *truth tables*.

# PRECISE MEANING: TRUTH TABLE

LaTeX

\wedge

\vee

\neg

English (or any natural human language) can be imprecise, so instead of using our “intuition” we **define** what “and” and “or” and “not” mean using **truth tables**.

and  $\wedge$

$p$	$q$	$p \wedge q$
1	1	1
1	0	0
0	1	0
0	0	0

or  $\vee$

$p$	$q$	$p \vee q$
1	1	1
1	0	1
0	1	1
0	0	0

not  $\neg$

$p$	$\neg p$
1	0
0	1

1 = True  
0 = False

define

# PRECISE MEANING: TRUTH TABLE

English (or any natural human language) can be imprecise, so instead of using our “*intuition*” we **define** what “*and*” and “*or*” and “*not*” mean using *truth tables*.

$p$	$q$	$p \wedge q$
1	1	1
1	0	0
0	1	0
0	0	0

$p$	$q$	$p \vee q$
1	1	1
1	0	1
0	1	1
0	0	0

$p$	$\neg p$
1	0
0	1

Teenager speech is more precise: Eg: “Maths is awesome — NOT”



# TRUTH TABLES FOR COMPOUND STATEMENTS

We can use truth tables to decide the truth values of more complicated statements, like  $\neg p \vee q$ :

$p$	$q$	$(\neg p) \vee q$
1	1	0
1	0	0
0	1	1
0	0	1

① ②



final true values.

~~$\neg (p \vee q)$~~

# TRUTH TABLES FOR COMPOUND STATEMENTS

We can use truth tables to decide the truth values of more complicated statements, like  $\neg p \vee q$ :

$p$	$q$	$\neg p \vee q$
1	1	1
1	0	0
0	1	1
0	0	1

↑

$p$	$q$	$\neg(p \vee q)$
1	1	0
1	0	0
0	1	0
0	0	1

↑

Note that this is different to saying  $\neg(p \vee q)$ , since the truth values are not the same

# YOUR TURN

$p$	$q$	$p \wedge q$
1	1	1
1	0	0
0	1	0
0	0	0

$p$	$q$	$p \vee q$
1	1	1
1	0	1
0	1	1
0	0	0

$p$	$\neg p$
1	0
0	1

Complete the truth tables for these statements:

$p$	$q$	$\neg (p \wedge q)$
1	1	0
1	0	1
0	1	1
0	0	1

②  
↑

①  
—

$p$	$q$	$\neg p$	$\vee$	$\neg q$
1	1	0	0	0
1	0	0	1	1
0	1	1	1	0
0	0	1	1	1

①  
↑

②  
↑

①  
—

$p$	$q$	$p \wedge q$	$\neg (p \wedge q)$
1	1	1	0
1	0	0	1
0	1	0	1
0	0	0	1

# LOGICALLY EQUIVALENT

When two (compound) statements have the same truth values we say they are logically equivalent.

$$\neg (p \wedge q)$$

IS  
LOG  
EQUIV.  
TO

$$\neg p \vee \neg q$$

# LOGICALLY EQUIVALENT

$$\neg q \vee \neg(\neg p)$$

When two (compound) statements have the same truth values we say they are *logically equivalent*.

Eg:

$p$	$q$	$p \vee \neg q$	$\neg(q \wedge \neg p)$
1	1	1 0	1 0 0
1	0	1 1	1 0 0
0	1	0 0	0 1 1
0	0	1 1	1 0 1
		(2) (1)	(3) (2) (1)

$p$	$\neg(\neg p)$
1	1
0	1

↑

↑

LOG EQUIV

$\neg p$

# IMPLIES

In mathematics and logic we have a very specific meaning for “ $p$  implies  $q$ ”, or “if  $p$  then  $q$ ”, notation  $p \rightarrow q$ .

We define it using the following table:

$p$	$q$	$p \rightarrow q$
1	1	1
1	0	0
0	1	1
0	0	1

# IMPLIES

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$p$	$q$	$p \rightarrow q$
1	1	1
1	0	0
0	1	1
0	0	1




You may think that in English, “if it is raining then I get wet” means that the rain caused me to get wet. But in mathematics if-then has the meaning defined above: if “I am wet” is true and “it is raining” is false, the implication is still true.

→ Canvas - link.

# YOUR TURN

Show that  $p \rightarrow q$  is logically equivalent to  $\neg p \vee q$ .

$p$	$q$	$p \rightarrow q$	$\neg p$	$\vee q$
1	1	1	0	1
1	0	0	0	0
0	1	1	1	1
0	0	1	1	1
			①	②





# TAUTOLOGY

A statement that is true for all truth value assignments is called a *tautology*.

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A statement that is true for all truth value assignments is called a *tautology*.

Eg:

$p$	$q$	$((p \rightarrow q) \wedge p) \rightarrow q$		
1	1	1	1	1
1	0	0	0	1
0	1	1	0	1
0	0	1	0	1

① ② ③

↑ Tautology.

# TAUTOLOGY

Eg:

$p$	$q$	$r$	$[(p \rightarrow q) \wedge (q \rightarrow r)] \rightarrow (p \rightarrow r)$				
1	1	1	1	1	1	1	1
1	1	0	1	0	0	1	0
1	0	1	0	0	1	1	1
1	0	0	0	0	1	1	0
0	1	1	1	1	1	1	1
0	1	0	1	0	1	1	1
0	0	1	1	1	1	1	1
0	0	0	1	1	1	1	1

3  
2

① ② ③ ④

$$\frac{(p \rightarrow q) \quad (q \rightarrow r)}{(p \rightarrow r)}$$

## YOUR TURN

Decide which of these are tautologies:

1.  $((p \rightarrow q) \wedge \neg q) \rightarrow \neg p$  *modus tollens.*
2.  $((p \rightarrow q) \wedge \neg\neg p) \rightarrow \neg q$

Draw TABLE!!

p	q	$((p \rightarrow q) \wedge \neg q)$			$\rightarrow$	$\neg p$
1	1	1	0	0	1	0
1	0	0	0	1	1	0
0	1	0	0	0	1	1
0	0	1	1	1	1	1

① ③ ② ⑤ ④

Tautology - all 1's.

Final Column

p	q	$((p \rightarrow q) \wedge \neg \neg p)$			$\rightarrow$	$\neg q$
1	1	1	1	0	0	0
1	0	0	0	1	1	1
0	1	0	0	1	1	0
0	0	0	0	1	1	1

①

④

same as p

③

②

⑥

⑤

Final Column

NOT A  
TAUTOLOGY.

## ANOTHER LOGICAL CONNECTIVE

Left right arrow.

Define  $\leftrightarrow$  ("if and only if") by the truth table:

$p$	$q$	$p \leftrightarrow q$
1	1	1
1	0	0
0	1	0
0	0	1

← Definition

## ANOTHER LOGICAL CONNECTIVE

Define  $\leftrightarrow$  (“if and only if”) by the truth table:

$p$	$q$	$p \leftrightarrow q$
1	1	1
1	0	0
0	1	0
0	0	1

$(p \rightarrow q) \wedge (q \rightarrow p)$ 

1
0
0
1

Ex: show that  $p \leftrightarrow q$  is logically equivalent to  $(p \rightarrow q) \wedge (q \rightarrow p)$

# YOUR TURN

Decide which of these are tautologies:

3.  $(p \rightarrow q) \leftrightarrow (q \rightarrow p)$

4.  $(p \rightarrow q) \leftrightarrow (\neg q \rightarrow \neg p)$

p	q	$(p \rightarrow q) \leftrightarrow (q \rightarrow p)$		
1	1	1	1	1
1	0	0	0	1
0	1	1	0	0
0	0	1	1	1

not a tautology

p	q	$(p \rightarrow q) \leftrightarrow \neg q \rightarrow \neg p$
		1
		1
		1
		1

YES  
TAUTOLOGY.



## ANOTHER WAY TO WRITE TAUTOLOGIES

In Humanities/Law you might see tautological statements written in this form. Some rules have names.

$$\begin{array}{c} p \rightarrow q \\ p \\ \hline q \end{array}$$

...

(Modus ponens)

$$\left( (p \rightarrow q) \wedge p \right) \rightarrow q$$

## ANOTHER WAY TO WRITE TAUTOLOGIES

In Humanities/Law you might see tautological statements written in this form. Some rules have names.

$$\frac{p \rightarrow q}{p}$$
$$\frac{p}{q}$$

(Modus ponens)

$$\frac{p \rightarrow q}{\neg q}$$
$$\frac{\neg q}{\neg p}$$

(Modus tollens)

$$\left( (p \rightarrow q) \wedge \neg q \right)$$

$$\rightarrow \neg p$$

## FROM WIKIPEDIA:

$M \rightarrow A$

If, I am an axe murderer, then I can use an axe.  
I cannot use an axe.  
Therefore, I am not an axe murderer.

Which style of argument is this? (Write it in symbols).

$\neg ( (M \rightarrow A) \wedge \neg A ) \rightarrow M$

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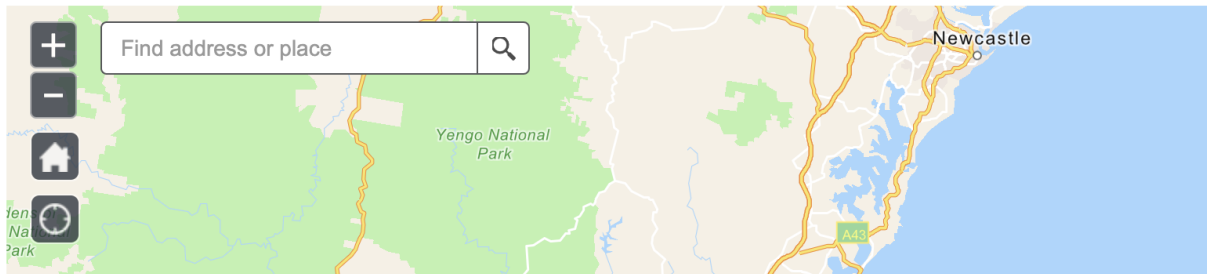
modus tollens

# OR, IF-THEN

## Metropolitan Sydney

If you live or work in the City of Sydney, Waverley, Randwick, Canada Bay, Inner West, Bayside, and Woollahra local government areas, you cannot travel outside metropolitan Sydney.

If you live or work in metropolitan Sydney, other than those local government areas, you may travel anywhere in NSW.



## CONTRADICTION: PREVIEW

Let  $F$  be a statement that is always false (has truth table 0, for example,  $F = q \wedge \neg q$ ).

Then the statement

$$(\neg p \rightarrow F) \rightarrow p$$

is a tautology. Check it:

$p$	$F$	$(\neg p \rightarrow F) \rightarrow p$	
1	0	0	1
0	0	1	1

①

②

③

Tautology.

$$p \wedge \neg p$$

## CONTRADICTION: PREVIEW

Let  $F$  be a statement that is always false (has truth table 0, for example,  $F = q \wedge \neg q$ ).

Then the statement

$$(\neg p \rightarrow F) \rightarrow p$$

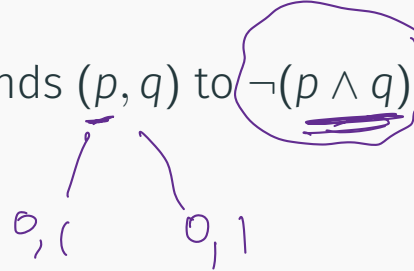
is a tautology. Check it:

$p$	$F$	
1		
0		

It says, if not  $p$  implies something that is false, then it must be  $p$  (is true). This argument form is known as *proof by contradiction*. We will study this more when we start *proofs*

A boolean function is a function from  $\{0, 1\}^n$  to  $\{0, 1\}$ . We will learn more about formal notation for functions in Week 3.

For example, the function  $f$  which sends  $(p, q)$  to  $\neg(p \wedge q)$  is a boolean function.

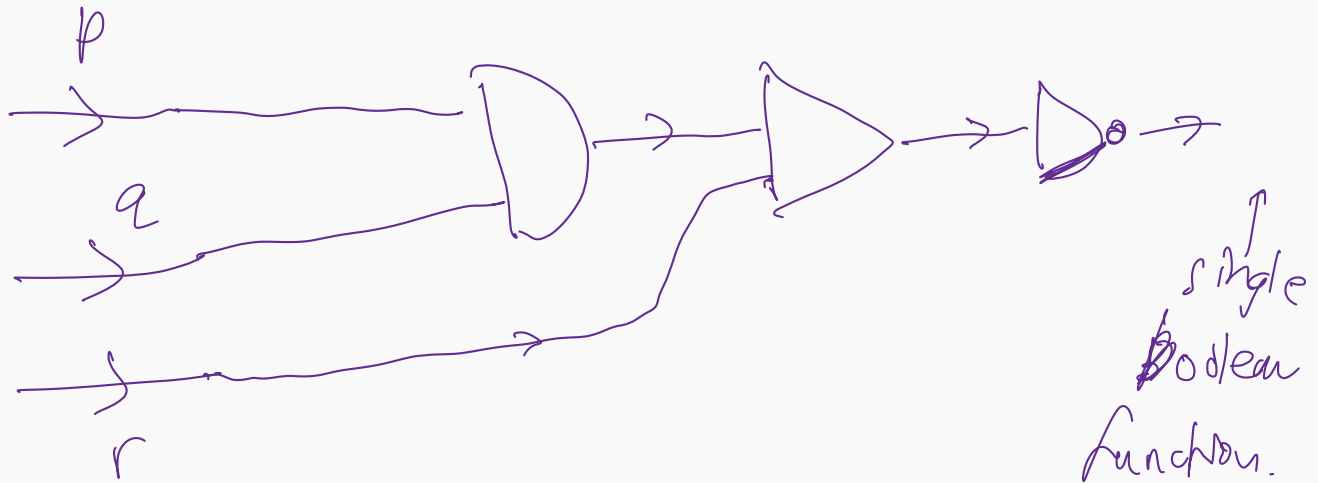


$(0, 0)$	$\rightarrow$	1
$(0, 1)$	$\rightarrow$	1
$(1, 0)$	$\rightarrow$	1
$(1, 1)$	$\rightarrow$	0

# LOGIC CIRCUITS

We can represent boolean functions as logic circuits which are theoretical models of a computer:

- input wires labeled by  $p, q, r, \dots$
- gates *AND, OR, NOT* ---
- single output wire, labeled by  $f(p, q, r, \dots)$ . *function*





# LOGIC CIRCUITS

We can represent boolean functions as logic circuits which are theoretical models of a computer:

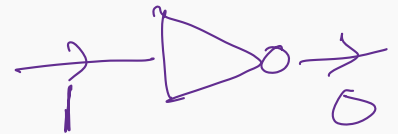
- input wires labeled by  $p, q, r, \dots$
- gates
- single output wire, labeled by  $f(p, q, r, \dots)$ .

Here are three kinds of gates used:

and : 

or : 

not: 



$p$	$q$	$p \wedge q$
1	1	1
1	0	0
0	1	0
0	0	0

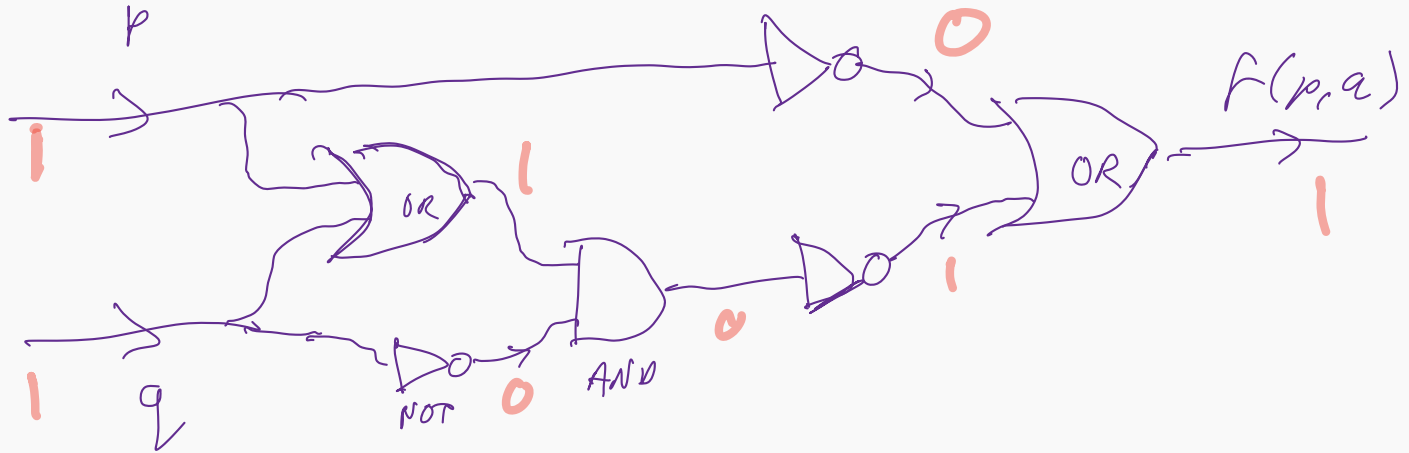
$p$	$q$	$p \vee q$
1	1	1
1	0	1
0	1	1
0	0	0

$p$	$\neg p$
1	0
0	1

# EXERCISE

Draw a logic circuit for the boolean function

$$f(p, q) = \neg((p \vee q) \wedge \neg q) \vee \neg p$$



# COMING UP

Next lecture (recorded, uploaded to Canvas by Tuesday 9am)

- quantified statements
- negation of quantified statements
- SAT and P=?NP

Please go through Lecture 2 recording with the blank slides, pause and do the problems as you watch.

Make sure you have seen all the content in lectures 1 and 2 before your tutorial class.

Note you are *not* expected to look at the Tutorial worksheet before your class, they are designed for you to work on together on-the-spot with your teammates during the class.

After your tutorial, do the Homework sheet by yourself (or with classmates if you can organise that).