

# 37181 DISCRETE MATHEMATICS

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Lecture 7: induction



- principle of mathematical induction

## HOW TO PROVE

### Lemma

For all  $n \in \mathbb{N}$ ,  $11^n - 4^n$  is divisible by 7.

$n=1$   
 $11-4=7$

$7 \mid 0$

$11^0 - 4^0 = 1 - 1 = 0$

$0 = 7 \cdot 0$

### Lemma

If  $A$  is a set of size  $n \in \mathbb{N}$ , then  $\mathcal{P}(A)$  has size  $2^n$ .

$n=1$   
 $A = \{a\}$   
 $\mathcal{P}(A) = \{\emptyset, \{a\}\}$   
 $|\mathcal{P}(A)| = 2^1$

$\mathcal{P}(\emptyset) = \{\emptyset\}$   
 $|\mathcal{P}(\emptyset)| = 1$

## Axiom (Principle of mathematical induction)

Let  $P(n)$  be a statement about natural numbers. Let  $s \in \mathbb{N}$ , eg.  
 $s = 0, 1$

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If

1.  $P(s)$  is true
2.  $P(k) \rightarrow P(k+1)$  is true



$$P(1)$$

$$P(1) \rightarrow P(2)$$

$$P(2) \rightarrow P(3)$$

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1.  $P(s)$  is true
2.  $P(k) \rightarrow P(k + 1)$  is true

then  $P(n)$  is true for all  $n \geq s$ .

$n \in \mathbb{N}$

# APPLICATION

## Lemma

For all  $n \in \mathbb{N}, n \geq 1$

$$\underline{1 + 2 + 3 + \dots + n} = \frac{n(n+1)}{2}$$

Proof Let  $P(n)$  be the statement:  
"  $1 + 2 + \dots + n = \frac{n(n+1)}{2}$  "

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$P(1)$  : LHS = 1 true.  
RHS =  $\frac{1(1+1)}{2} = \frac{1 \cdot 2}{2} = 1$   
LHS = RHS so  $P(1)$  is true

Assume  $p(k)$  is true for some  $k \geq 1$ .  
 $\rightarrow (1 + \dots + k = \frac{k(k+1)}{2})$  \*

To show:  $p(k+1)$ :

$$\text{LHS} = 1 + 2 + \underline{3} + 4 + \dots + \underline{k} + \underline{(k+1)}$$

$$= \frac{k(k+1)}{2}$$

$$+ (k+1)$$

by Inductive Assumption \*

$$= (k+1) \left( \frac{k}{2} + 1 \right)$$

$$= (k+1) \left( \frac{k}{2} + \frac{2}{2} \right)$$

$$= \frac{(k+1)(k+2)}{2} = \text{RHS.}$$

thus  $p(k) \rightarrow p(k+1)$ .

Thus by P.O.I.,  $p(n)$  is true for all  $n \geq 1$ .

# APPLICATION

## Lemma

For all  $n \in \mathbb{N}, n \geq 1$

$$1 + 2 + 3 + \cdots + n = \frac{n(n+1)}{2}$$

## Proof.

Let  $P(n)$  be the statement that  $1 + 2 + 3 + \cdots + n = \frac{n(n+1)}{2}$ .

Don't  
write  $P(n) =$

# APPLICATION

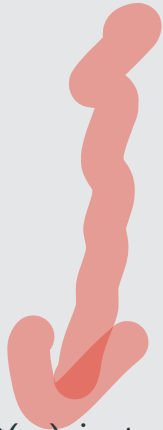
## Lemma

For all  $n \in \mathbb{N}, n \geq 1$

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## Proof.

Let  $P(n)$  be the statement that  $1 + 2 + 3 + \cdots + n = \frac{n(n+1)}{2}$ .



Thus by PMI  $P(n)$  is true for all  $n \geq 1$ .



# APPLICATION

## Lemma

For all  $n \in \mathbb{N}, n \geq 1$

$$1^2 + 2^2 + 3^2 + \dots + n^2 = \frac{n(n+1)(2n+1)}{6}$$

Let  $P(n)$  be statement  
" $1^2 + 2^2 + \dots + n^2 = \frac{n(n+1)(2n+1)}{6}$ "

---

$$P(1): \text{LHS} = 1^2 = 1$$
$$\text{RHS} = \frac{1(1+1)(2+1)}{6} = \frac{1 \cdot 2 \cdot 3}{6} = 1$$

So  $P(1)$  true

Assume  $P(k)$  (\*) is true for some  $k \geq 1$

To show:  $P(k+1)$

$$\text{LHS} = 1^2 + 2^2 + \dots + k^2 + (k+1)^2$$

$$= \frac{k(k+1)(2k+1)}{6} + (k+1)^2$$

by Ind assumption (\*)

$$= (k+1) \left( \frac{k(2k+1)}{6} + (k+1) \right)$$

$$= (k+1) \left( \frac{2k^2 + k + 6k + 6}{6} \right)$$

$$= (k+1) \frac{(k+2)(2k+3)}{6}$$

$$= \text{RHS}$$

$$\text{So } P(k) \rightarrow P(k+1)$$

So by PMI,  $P(n)$  is true for all  $n \geq 1$ .

# APPLICATION

## Lemma

For all  $n \in \mathbb{N}, n \geq 1$

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Let  $P(n)$  be the statement that



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## Lemma

For all  $n \in \mathbb{N}, n \geq 1$

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Thus by PMI  $P(n)$  is true for all  $n \geq 1$ .



incl 0.

## Lemma

For all  $n \in \mathbb{N}$ ,  $11^n - 4^n$  is divisible by 7.Let  $P(n) : \quad \neg \mid (11^n - 4^n)$  " "

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$$P(0) : n=0, \quad 11^n - 4^n = 11^0 - 4^0 = 1 - 1 = 0$$

and  $\neg \mid 0$  since  $0 = 7 \cdot 0$

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Assume  $P(k)$  true.
 $\exists s \in \mathbb{Z}$   
 so that
To show:  $P(k+1) :$ 

$$\Rightarrow 11^k - 4^k$$

$$11$$

$$7 \cdot s$$

$$\begin{aligned}
& 11^{k+1} - 4^{k+1} \\
&= 11 \cdot \underline{11^k} - 4 \cdot \underline{4^k} \\
&= (7 + 4) 11^k - 4 - 4^k \\
&= 7 \cdot 11^k + 4(11^k - 4^k)
\end{aligned}$$

divisible by 7

$$= 7 \cdot 11^k + 4 \cdot 7s$$

$$= 7(11^k + 4s)$$

$$\text{So } 7 \mid 11^{k+1} + 4^{k+1}$$

So  $P(k+1)$  is true

So  $P(k) \rightarrow P(k+1)$

So by PMI  $P(n)$  true  $\forall n \geq 0$ .

**Lemma**

*For all  $n \in \mathbb{N}$ ,  $11^n - 4^n$  is divisible by 7.*

**Proof.**

Let  $P(n)$  be the statement that

**Lemma**

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**Proof.**

Let  $P(n)$  be the statement that



Thus by PMI  $P(n)$  is true for all  $n \geq 0$ .



## Lemma

If  $A$  is a set of size  $n \in \mathbb{N}$ , then  $\mathcal{P}(A)$  has size  $2^n$ .

## Proof.

Let  $P(n)$  be the statement that if  $|A| = n$   
then  $|\mathcal{P}(A)| = 2^n$ .

---

$p(0)$  : The only set of size 0  
is  $\emptyset$ , and  $\mathcal{P}(\emptyset) = \{\emptyset\}$   
so  $|\mathcal{P}(\emptyset)| = 1 = 2^0$  so  $p(0)$   
is true

~~Thus by PMI  $P(n)$  is true for all  $n \geq 0$ .~~

□

Assume  $P(k)$  is true.

To show:  $P(k+1)$ :

Let  $A$  be some set with  $|A| = k+1$ .

$$A = \{a_1, a_2, a_3, \dots, a_{k+1}\}$$

Consider all the subsets  $S$  that do not contain  $a_1$ .

Put all the subsets that do not contain  $a_1$ .

Put all subsets that do contain  $a_1$ .

$\{a_1\}$

$\{a_1, a_2\}$

$2^k$

$\emptyset$

$\{a_2\}$

$\{a_2, a_3\}$

$2^k$

same number

$\{a_2, a_3, \dots, a_{k+1}\}$

$\{a_1, a_2, \dots, a_{k+1}\}$

$a_2, a_3, \dots, a_{k+1}$

Either a subset of  $A$  contains  $a_1$ , or not

all subsets  
of  $\{ \underbrace{a_2, a_3, \dots, a_{k+1}}_{\text{only } k \text{ elements}} \}$

by inductive assumption,  
there are  $2^k$  subsets  
on the left side,

and  $2^k$  on right  
side.

$$\begin{aligned} 2^k + 2^k &= 2 \cdot 2^k \\ &= 2^{k+1} \end{aligned}$$

— Please use

## Lemma

For all  $n \in \mathbb{N}$ , if  $n \geq \square$  then (some statement).

## Proof.

Let  $P(n)$  be the statement

Then  $P(\square)$  is true since

Assume  $P(k)$  for  $k \geq \square$ . Then

$\therefore P(k+1)$  is true

Thus by PMI  $P(n)$  is true for all  $n \geq \square$ .

□

$P(k):$   
~ ~ ~

# STRONGER VERSION (OR IS IT?)

PMI is equivalent to the following: Let  $s \in \mathbb{N}$ .

If

- $P(s)$  is true and
- if for all  $s \leq i \leq n$   $P(i)$  is true, then  $P(n + 1)$  is true,

## STRONGER VERSION (OR IS IT?)

PMI is equivalent to the following: Let  $s \in \mathbb{N}$ .

If

- $P(s)$  is true and
- if *for all*  $s \leq i \leq n$   $P(i)$  is true, then  $P(n + 1)$  is true,

then  $P(n)$  is true for all  $n \in \mathbb{Z}, n \geq s$ .

## Lemma

For all  $n \in \mathbb{N}, n > 1$  if  $n$  is not prime then some prime number  $p$  divides  $n$ .

## Proof.

(Strong PMI)

Let  $P(n)$  statement that if  $n$  is not prime then  $\exists p$  prime so that  $p|n$ .

  $P(2)$ : 2 is a prime so  $P(2)$  is true

Assume  $P(2), P(3), \dots, P(n)$  is true.  
 $n \geq 2$

To show:  $P(n+1)$ .

If  $n+1$  is prime,  $P(n+1)$  is true

Else,  $n+1$  is not prime, so  $\exists a, b \in \mathbb{N}$   
so that  $(n+1) = a \cdot b$   $a, b \neq 1$

$a, b \geq 2$

$$2 \leq a \leq n$$

So by Inductive assumption,  
 $P(a)$  is true

so  $\exists p$  prime so that  $p \mid a$   
 $a = p \cdot s$   $s \in \mathbb{N}$

$$\begin{aligned} n+1 &= a \cdot b \\ &= p \cdot s \cdot b \end{aligned}$$

So  $p \mid n+1$ .

So  $P(n+1)$  is true.

so by SPMI  $P(n)$  is true  $\forall n \geq 2$

### Lemma

*For all  $n \in \mathbb{N}, n > 1$  if  $n$  is not prime then some prime number  $p$  divides  $n$ .*

### Proof.

Let  $P(n)$  be the statement that either  $n$  is prime or some prime divides  $n$ .



$$n = 0 \quad \begin{array}{l} \text{LHS} = 0! = 1 \\ \text{RHS} = 2^{-1} = \frac{1}{2} \end{array}$$

## Lemma

For all  $n \in \mathbb{N}$ ,  $n! \geq 2^{n-1}$

## Proof.

Let  $P(n)$  be the statement that

$$n! \geq 2^{n-1}$$

~~$P(0)$  is true since  $0! = 1$  and  $2^{-1} = \frac{1}{2}$~~

~~$P(1)$  is true  $1! \geq 2^{1-1} = 1$~~

Assume  $P(k)$  ~~\*~~ is true,  $k \geq 0$

To show  $P(k+1)$ :  $(k+1)! = (k+1)k!$

(start at 0)

$$\geq (k+1) 2^{k-1}$$

by Inductive  
Assumption \*

$$\begin{aligned} &\geq (1+1) 2^{k-1} \quad \text{since } k \geq 1 \\ &= 2 \cdot 2^{k-1} = 2^k \\ &= 2^{(k+1)-1} \end{aligned}$$

So  $p(k+1)$   
is true.

---

So by PMI,  $p(n)$  is true  
for all  $n \geq 1$   
and we also showed it is  
true for  $n = 0$ ,

So  $p(n)$  is true for  $n \geq 0$ .

---

$$n=0 \checkmark$$

$$n=1$$

$$1! = 1$$

Consider the statement: For all  $n \in \mathbb{N}$ ,  $n! \geq 3^{n-1}$ .

$$3 = 1$$

Is this true? Is it true for all  $n \geq ??$

**Proof.**

Let  $P(n)$  be the statement that

$$n=2$$

$$2! = 2$$

$$3^{2-1} = 3$$



$$n=3$$

$$3! = 6$$

$$3^{3-1} = 9$$

□



$$4! = 24$$

$$3^3 = 27$$



$$n \cdot (n-1) \cdot (n-2) \cdots 3 \cdot 2 \cdot 1$$

$$3 \cdot 3 \cdot 3 \cdots 3 \cdot 3 \cdot 1$$

(start at ??)

## Lemma

All horses are black.



Proof:

$n=0$

I have no horses

Which is true:

They are all  
black



If I have  $n \in \mathbb{N}$   
horses, then they are  
all black.

NOT

$\exists$  one horse  
that is  
not black.



So  $P(0)$  is true.

---

Assume (for induction) that for any collection of  $k \geq 0$  horses,  $k \in \mathbb{N}$ , they are all black.

---

Now, consider some collection of  $k+1$  horses.



Remove one horse

→ you get  $k$  horses left.

by Ind. Assumption, these horses are all black.

---

---

Put that horse back  
and remove another horse

---

leaves  $k$  horses  
 $\Rightarrow$  ~~now~~ that horse must be  
black horse.

---

by PMI, all horses are black.

**Lemma**

*All horses are black.*

**Proof.**

Let  $P(n)$  be the statement that



Next lecture:

- application of induction: correctness of computer code/algorithms
- WOP and PMI  
~~WOP~~ ~~PMI~~

Important to gets lots of practice doing proofs by induction.

