

DISCRETE MATH 37181 TUTORIAL WORKSHEET 4

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INSTRUCTIONS. Complete these problems in groups of 3-4 at the whiteboard. Partial solutions at the end of the PDF.

1. Prove (use induction) that for all $n \in \mathbb{N}_+$

$$2 + 4 + 6 + 8 + \cdots + (2n) = n^2 + n$$

2. Prove (use induction) that for all $n \in \mathbb{N}_+$

$$2 + 7 + 12 + 17 + 22 + \cdots + (5n - 3) = \frac{n(5n - 1)}{2}$$

3. Prove that 6 divides $n^3 + 5n$ for all $n \in \mathbb{N}$.

4. Let $P(n)$ be the statement that $n^2 + 5n + 1$ is even.

(a) Prove that $P(n)$ implies $P(n + 1)$ for any $n > 0$.

(b) For which values of n is $P(n)$ actually true?

(c) What is the *moral* of this exercise?

5. Let $c \in \mathbb{R}_+ = \{x \in \mathbb{R} \mid x > 0\}$. Prove that $\exists k \in \mathbb{N}$ such that for all $n \geq k$, $n! > c^n$. ^{1 2}

6. Prove or disprove:

$$\forall n \in \mathbb{N}, \quad n^3 + 4n \equiv 0 \pmod{5}$$

7. Prove or disprove: For all $n \in \mathbb{N}_+$,

$$\sum_{i=1}^n \frac{1}{i(i+1)} = \frac{n}{n+1}.$$

8. Find an expression (which does not use $+\cdots+$ or $\sum$) for the function **unknown** computed by the following code:

```
int unknown(int n)
{
  if n=1 return 1
  else return (unknown(n-1)+(2*n-1)^2)
}
```

Date: Week 4 workshop (Wednesday 16, Thursday 17, Friday 18 March).

¹this means that *eventually* the function $n!$ will *dominate* or be bigger than c^n .

²Hint: If there are too many abstract things here, try doing the question for some *fixed* value of $c \in \mathbb{R}_+$ first. Eg $c = 5$. Once you understand how to do it for $c = 5$, try repeating it for a general positive real number c instead.

9. Prove that PMI implies WOP. Start with:

Suppose (for contradiction) that WOP is false, so there is some non-empty set S which does not have a first element.

Let $P(n)$ be the statement that “for all $i \in \mathbb{N}, i \leq n, i \notin S$ ”.

10. Complete this definition: A *loop invariant* is ...

11. Consider this pseudocode.

```
function mycode(int m, int n)
{
  while (m >= 0 and m <= 100)
    m := m+1
    n := n-1
  end while
  return n
}
```

- Show that “ $m + n = 50$ ” is a loop invariant for the while loop.
- Show that “ $m + n$ is even” is a loop invariant for the while loop.
- Show that “ $m + n$ is odd” is a loop invariant for the while loop.
- Why does it terminate? What is the output?

m	n	output
100	0	
-10	100	
50	0	
0	0	

(You can use this table to try out the algorithm on different inputs, then go back and answer part (d).)

12. Recall problem 3 (vi) from worksheet 2. At that time you had just seen the WOP but not really used it. Now the following statement should be straightforward to prove using WOP³.

For all $n \in \mathbb{Z}$, if $n > 1$ then n is divisible by some prime number p .

Hint: Given $n > 1$, let $D = \{s \in \mathbb{Z} \mid (s|n) \wedge (s > 1)\}$ be the set of integers that divide n . Explain why D is non-empty, and also why $D \subseteq \mathbb{N}$. Then show that the first element of D is a prime which divides n .

13. Prove that if $x > 0$ is any fixed real number, then

$$(1 + x)^n > 1 + nx$$

for all $n \in \mathbb{N}, n \geq 2$.

14. Prove or disprove: for all $n \in \mathbb{N}$, $(3n + 1)7^n - 1$ is divisible by 9.

³or induction, which we did in the lecture, but for this question try to prove it using WOP.

Brief solutions:

1. Let $P(n)$ be the statement that

$$2 + 4 + 6 + 8 + \cdots + (2n) = n^2 + n.$$

Then $P(1)$ is true since LHS=2 and RHS= $1 + 1 = 2$.

Assume $P(k)$ is true. Then $P(k+1)$: LHS=

$$2 + 4 + 6 + 8 + \cdots + (2k) + (2k+2) = k^2 + k + (2k+2)$$

using the inductive assumption

$$\begin{aligned} &= k^2 + 3k + 2 = (k+1)(k+2) = (k+1)((k+1)+1) = (k+1)^2 + (k+1) \\ &= \text{RHS} \end{aligned}$$

Then by PMI $P(n)$ is true for all $n \geq 1$.

2. Let $P(n)$ be the statement that

$$2 + 7 + 12 + 17 + 22 + \cdots + (5n-3) = \frac{n(5n-1)}{2}.$$

Then $P(1)$ is true since LHS=2 and RHS= $\frac{1(4)}{2} = 2$.

Assume $P(k)$ is true. Then $P(k+1)$:

$$\text{LHS} = 2 + 7 + 12 + 17 + 22 + \cdots + (5k-3) + (5k+5-3) = \frac{k(5k-1)}{2} + (5k+2)$$

using the inductive assumption

$$= \frac{k(5k-1)}{2} + \frac{2(5k+2)}{2} = \frac{5k^2 - k + 10k + 4}{2} = \frac{5k^2 + 9k + 4}{2}$$

(secretly I will work out the RHS, then make them match up
(and write it nicely at the end))

$$\text{RHS} = \frac{(k+1)(5k+5-1)}{2} = \frac{(k+1)(5k+4)}{2} = \uparrow$$

Then by PMI $P(n)$ is true for all $n \geq 1$.

3. Let $P(n)$ be the statement that 6 divides $n^3 + 5n$, then $P(1)$ is true since $1^3 + 5 \cdot 1 = 6$. Since we need to prove for all $n \in \mathbb{N}$, we must start at $n = 0$: $P(0)$ is true since 6 divides $0 = 0^3 + 5 \cdot 0$.

Assume for induction that $P(k)$ is true for $k \geq 0$, which means $k^3 + 5k = 6s$ for some $s \in \mathbb{Z}$. Then

$$(k+1)^3 + 5(k+1) = k^3 + 3k^2 + 3k + 1 + 5k + 5 = k^3 + 3k^2 + 8k + 6$$

Now use the inductive assumption

$$= k^3 + 5k + 3k + 3k^2 + 6 = 6s + 3k(k+1) + 6$$

so if we can show $3k(k+1)$ is a multiple of 6 we are done. It is because $k(k+1)$ is even, since either k is even or if not $k+1$ is even. Thus $P(k)$ implies $P(k+1)$.

Then by PMI $P(n)$ is true for all $n \geq 0$.

4. (a) Assuming $P(n)$ we have $n^2 + 5n + 1 = 2s$ for some $s \in \mathbb{Z}$. Then

$$P(n+1) : (n+1)^2 + 5(n+1) + 1 = n^2 + 2n + 1 + 5n + 5 + 1 = n^2 + 5n + 1 + 2n + 6 = 2s + 2n + 6$$

is even, so $P(n) \rightarrow P(n+1)$.

(b) None.

(c) Moral: a proof by induction needs *both* $P(1)$ true (or some starting number) and $P(n) \rightarrow P(n+1)$ to work.

5. Let $P(n)$ be the statement that $n! > c^n$.

Assume $P(k)$, then $P(k+1) : (k+1)! = k!(k+1) > c^k(k+1) > c^k \cdot c$ if we assume that $k+1 > c$ and $c^k \cdot c = c^{k+1}$, so $P(k) \rightarrow P(k+1)$ assuming we choose $k \geq c$.

Base step: we need to find a value of k so that $P(k)$ is actually true (remember the moral learned from the last exercise!).

$$\begin{array}{cccccccccccc} k! & = & k & (k-1) & (k-2) & \dots & ([c]+1) & [c] & ([c]-1) & \dots & 3 & 2 & 1 \\ c^k & = & c & c & c & \dots & c & c & c & \dots & c & c & c \end{array}$$

Using this picture, I claim that if I choose $k = \lceil c \rceil^2$ (rounding up the value of $c \in \mathbb{R}^+$ to the nearest whole number) then $(\lceil c \rceil^2)! > c^{\lceil c \rceil^2}$.

To prove this, I need another picture:

$$\begin{array}{cccccccccccc} k! & = & (\lceil c \rceil^2) & (\lceil c \rceil^2 - 1) & \dots & (\lceil c \rceil + 1) & \lceil c \rceil & (\lceil c \rceil - 1) & \dots & 3 & 2 & 1 \\ c^k & = & c & c & \dots & c & c & c & \dots & c & c & c \end{array}$$

(can you complete this proof?)

Then by PMI $P(n)$ is true for all $n \geq \lceil c \rceil^2$.

6. False: true for 0,1 but for $n = 2$ it is false.

7. This is true, and since its a statement about all n we use PMI.

Let $P(n)$ be the statement that $\sum_{i=1}^n \frac{1}{i(i+1)} = \frac{n}{n+1}$. Then $P(1)$ is true since LHS = $\frac{1}{1 \cdot 2}$ and RHS = $\frac{1}{1+1} = \frac{1}{2}$.

Assume $P(k)$, then consider $P(k+1)$:

$$\begin{aligned} \sum_{i=1}^{k+1} \frac{1}{i(i+1)} &= \sum_{i=1}^k \frac{1}{i(i+1)} + \frac{1}{(k+1)(k+2)} \\ &= \frac{k}{k+1} + \frac{1}{(k+1)(k+2)} \\ &= \frac{k(k+2)}{(k+1)(k+2)} + \frac{1}{(k+1)(k+2)} \\ &= \frac{k^2+2k+1}{(k+1)(k+2)} \\ &= \frac{(k+1)(k+1)}{(k+1)(k+2)} \\ &= \frac{(k+1)}{(k+2)} \end{aligned}$$

so $P(k+1)$ is true.

So by PMI since $P(1)$ is true and $P(k) \rightarrow P(k+1)$ then $P(n)$ is true for all $n \in \mathbb{N}_+$.

8. $u(1) = 1, u(n) = u(n-1) + (2n-1)^2$, so

$$u(n) = 1 + 3^2 + 5^2 + 7^2 + \cdots + (2n-1)^2.$$

If we want a formula that doesn't involve a sum, guess a formula for $u(n)$ by computing small values, then use induction to prove it.

Solution:

$$1 + 3^2 + \cdots + (2n-1)^2 = \frac{n(2n+1)(2n-1)}{3}$$

9. Proof: Suppose (for contradiction) that WOP is false, so there is some non-empty set S which does not have a first element.

Let $P(n)$ be the statement that “for all $i \in \mathbb{N}, i \leq n, i \notin S$ ”.

Then $P(0)$ is true because if 0 belongs to S , it would have to be the first element, and so it (and all $i \leq 0$, which is just $i = 0$) is not in S .

If $P(k)$ is true then $0, 1, 2, \dots, k \notin S$. If $k+1$ is in S then it would have to be the (recall we proved first elements, if they exist, must be unique) first element, but S does not have one, so $k+1$ is not in S , and so $P(k+1)$ is true.

Then by PMI $P(n)$ is true for every $n \in \mathbb{N}$ which means that S is empty, which is a contradiction since we started by saying S is non-empty.

10. A *loop invariant* is a statement that if true before one iteration of the loop, remains true after one iteration.

11. (a) If $m + n = 50$, then in one iteration $m + n \rightarrow m + 1 + n - 1 = m + n = 50$ still.

- (b) It terminates because each iteration increases the value of m , and if m exceeds 100 then it will stop. If m starts off negative or greater than 100 then it will not enter the loop at all at stop straight away.

Output: If $m < 0$ or $m > 100$ then it simply outputs the value n entered by the user.

If $0 \leq m \leq 100$ then we subtract stuff from n . Let's use the loop invariant to help us. Call m', n' the new values of the variables at the end of the loop.

At the end, $m' = 101$ (if $m = 100$ we perform one more iteration of the loop). We know that the sum $m + n$ remains invariant during the loop, so if originally $m + n = p$ then at the end $101 + n' = p$, so $n' = p - 101$. So I claim the output is $n + m - 101$.

m	n	output
100	0	-1
-10	100	100
50	0	-51
0	0	-101

13. Let $P(n)$ be the statement that $(1+x)^n > 1+nx$ where x is a fixed positive real number.

Then $P(2)$ is true since $(1+x)^2 = 1+2x+x^2 > 1+2x$ since $x^2 > 0$.

Assume $P(k)$ is true for $k \geq 2$. Then $P(k+1)$:

$$(1+x)^{k+1} = (1+x)^k(1+x) > (1+kx)(1+x) = 1+kx+x+kx^2 = 1+(k+1)x = kx^2 > 1+(k+1)x$$

since $kx^2 > 0$ so $P(k) \rightarrow P(k+1)$.

Then by PMI true for all $n \geq 2$.

14. Hint: You may need to prove another fact first: $7^m + 2$ is divisible by 3.