

31268 Web Systems

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Week 6: Operating Systems 3

Part 1 - Process Management

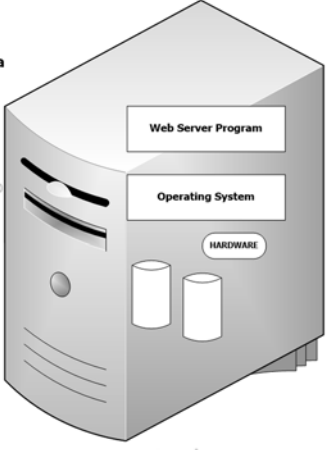
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Recap: The web...

A bunch of computers and a network of networks...

BROWSER

... and a whacking **big** computer running the web site program on an operating system running on **HARDWARE**



Web Server Program

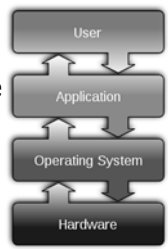
Operating System

HARDWARE

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Recap: Operating System

- Manages your computer
- Runs programs
- Interface between user and hardware
- Provides services to programs & users
- Protects users and programs from each other....



User

Application

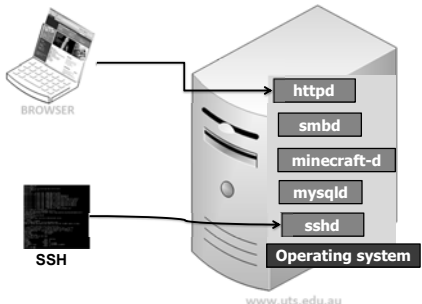
Operating System

Hardware

wikipedia 3

Web Example

Our Web Servers run various programs to provide services



BROWSER

SSH

httpd

smbd

minecraft-d

mysqld

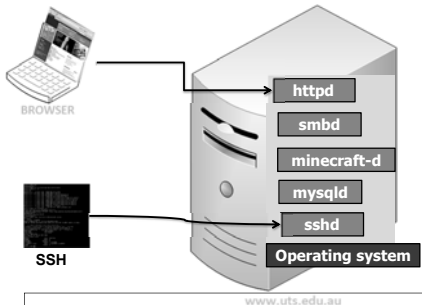
sshd

Operating system

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Web Example

Our Web Servers run various programs to provide services



BROWSER

SSH

httpd

smbd

minecraft-d

mysqld

sshd

Operating system

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- Today we will look at how O/S runs programs

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Operating Systems - Outline

- *Operating Systems*
- *file systems and file manipulation*
 - Disk physical structure
 - Disk logical structures
 - File allocation methods
 - Unix: files, directories and redirection
- *Scripting and regular expressions*
 - Unix and the Command-Line interface
 - Variables and bash scripts
- **processes, threads, piping and redirection**
 - Programs vs processes, IPC, deadlocks
- *memory and process management*
 - Physical, Logical & Virtual memory

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Programs vs Processes



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- Processes are programs in execution.
(also called "tasks")
- Every time a program is **run**, a new process is created.
- In addition there are "system processes", "services" or "background tasks"
 - Eg mouse, antivirus, ...



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Process Management

- The role of the computer is to run our programs
- Process management is how the OS handles this step
 - starting processes
 - managing running processes
 - performing interprocess communication
 - terminating processes

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Managing processes



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- Modern operating systems can **appear** to run multiple programs concurrently
- Run one program at a time
 - single tasking – *e.g. traditional embedded systems*
 - batch processing – *e.g. ancient mainframe*



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Managing processes



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- Modern operating systems can **appear** to run multiple programs concurrently
- Run one program at a time
 - single tasking – *e.g. traditional embedded systems*
 - batch processing – *e.g. ancient mainframe*
- Concurrent processing
 - Multiprogramming – *e.g. windows 3*
 - Multitasking – *all modern operating systems*
 - Multithreading – *most modern operating systems pentium 4+*
 - multiprocessing – *multicore CPUs, dual/quad CPU*



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Process States

Processes can be in various "states"



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Memory



wikipedia

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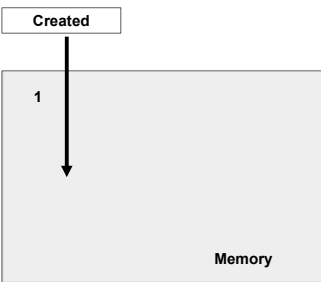
Process States

Processes can be in various "states"



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1. **Created** (ie: loaded into memory)



Created

1

Memory



wikipedia

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Process States

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Processes can be in various "states"

1. **Created** (ie: loaded into memory)
2. **Waiting** to be run by CPU

Created

1

2

Waiting

Memory

```
graph TD; Created[Created] -- 1 --> Waiting((Waiting)); Waiting -- 2 --> Waiting;
```

→

wikipedia

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Process States

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Processes can be in various "states"

1. **Created** (ie: loaded into memory)
2. **Waiting** to be run by CPU
3. O/S **runs** the process

Created

1

2

Waiting

3

Running

Memory

```
graph TD; Created[Created] -- 1 --> Waiting((Waiting)); Waiting -- 2 --> Waiting; Waiting -- 3 --> Running((Running)); Running -- 3 --> Running;
```

→

wikipedia

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Process States

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Processes can be in various "states"

1. **Created** (ie: loaded into memory)
2. **Waiting** to be run by CPU
3. O/S **runs** the process
4. If process needs a resource, "**blocked**" (ie: waiting) until it gets the resource.

Created

1

2

Waiting

3

Running

4

Blocked

Memory

```
graph TD; Created[Created] -- 1 --> Waiting((Waiting)); Waiting -- 2 --> Waiting; Waiting -- 3 --> Running((Running)); Running -- 4 --> Blocked((Blocked)); Blocked -- 4 --> Blocked;
```

→

wikipedia

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Process States

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Processes can be in various "states"

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5. It then waits again until the OS re-starts the process.

Created

1

2

Waiting

3

Running

4

Blocked

5

Memory

```
graph TD; Created[Created] -- 1 --> Waiting((Waiting)); Waiting -- 2 --> Waiting; Waiting -- 3 --> Running((Running)); Running -- 4 --> Blocked((Blocked)); Blocked -- 5 --> Waiting;
```

→

wikipedia

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Process States

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Created

1

2

Waiting

3

Running

4

Blocked

5

Memory

```
graph TD; Created[Created] -- 1 --> Waiting((Waiting)); Waiting -- 2 --> Waiting; Waiting -- 3 --> Running((Running)); Running -- 4 --> Blocked((Blocked)); Blocked -- 5 --> Waiting;
```

→

wikipedia

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Keep looping until finished

Process States

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Processes can be in various "states"

1. **Created** (ie: loaded into memory)
2. **Waiting** to be run by CPU
3. O/S **runs** the process
4. If process needs a resource, "**blocked**" (ie: waiting) until it gets the resource.
5. It then waits again until the OS re-starts the process.
6. When finished, O/S **stops** the process

Created

1

2

Waiting

3

Running

4

Blocked

5

6

Stopped

Memory

```
graph TD; Created[Created] -- 1 --> Waiting((Waiting)); Waiting -- 2 --> Waiting; Waiting -- 3 --> Running((Running)); Running -- 4 --> Blocked((Blocked)); Blocked -- 5 --> Waiting; Waiting -- 6 --> Stopped[Stopped];
```

→

wikipedia

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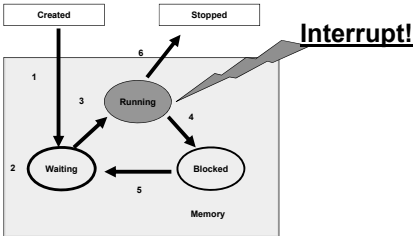
Interrupts

- The O/S runs processes continuously until interrupted
- Interruptions come from
 - Hardware
 - E.g. Ctrl-Alt-Delete keystroke, move mouse, network
 - Programs
 - E.g. runtime error, pause for input, wait for block of data to load from disk

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Interrupts

- O/S will force current running process into "blocked" state & runs a special "interrupt handler"
- Often it will pass control to a different process while waiting .



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Process Scheduling



- How does O/S decide when to run each process on which CPU?
→ This is called "**Scheduling**".
- Various types of scheduling exist but most use a "Queue"



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Process Scheduling



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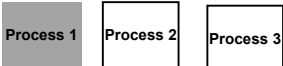


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Process Scheduling



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Process Scheduling

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Process 1

Process 2

Process 3

Process 4



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Process Scheduling

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Process 1

Process 2

Process 3

Process 4



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Process Scheduling

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Process 1

Process 2

Process 3

Process 4



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Process Scheduling

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Process 1

Process 2

Process 3

Process 4



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Process Scheduling

- How does O/S decide when to run each process on which CPU?
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Running process

Process 1

Process 2

Process 3

Process 4



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Process Scheduling

- How does O/S decide when to run each process on which CPU?
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Running process

Process 1

Process 2

Process 3

Process 4

Waiting Processes

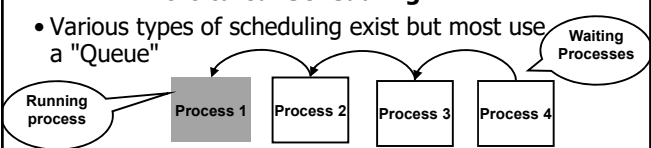


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Process Scheduling

- How does O/S decide when to run each process on which CPU?
→ This is called **"Scheduling"**.
- Various types of scheduling exist but most use a "Queue"



- Different types of algorithms to decide which process runs next eg:
 - First-in first-out
 - Pre-emptive
 - Round-robin
 - ... + heaps more!



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Concurrent programming

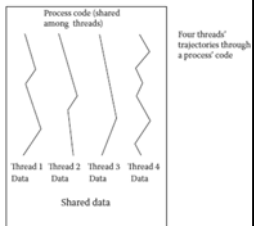
- Various O/S uses different techniques to run processes:
 - **Multiprogramming** – "cooperative"
 - when process waits for I/O, run next process
 - Can get "hangs" if a process does no I/O or doesn't relinquish the CPU
 - **Multitasking** - "pre-emptive"
 - Set a timer on all processes – "fair share" CPU time. When timer stops, switch to next process.
 - Can have priorities eg – low, normal, high, real-time
 - Choice of Queuing algorithm decides which process runs next

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Concurrent programming -2

- Various O/S uses different techniques to run processes:
 - **Multithreading**
 - Multiple instances of the same code eg same process in memory, but 2 or more execute paths
 - Eg powerpoint – might have 1 thread per presentation
 - Allows O/S to run a process on more than 1 CPU (e.g. 1 per concurrent thread)
 - Very efficient use of resources eg memory
 - **Multiprocessing**
 - Multitasking spread amongst 2 or more CPU's



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Inter Process Communication

- Processes often need to communicate with each other. This is called **IPC**.
- Operating Systems allow communication between processes via...
 - File sharing
 - Shared memory
 - Signals
 - Sockets
 - Messages
 - Pipes
 - Semaphores
 - Locks

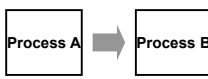


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IPC one-way example: Pipes

- One way communication.



- Process **A** sends data to process B. Process B accepts the data.
- EXAMPLE: Bash "pipe" operator
→ output of one command (process) becomes the input of another command (process).

Is | more

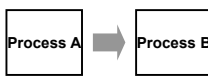


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IPC one-way example: Pipes

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Is → more

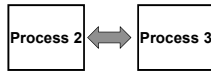


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Two-way IPC



- Processes can communicate in both directions
 - shared memory (bit of memory which acts as like a file)
 - Named pipe (special 2 way pipe, acts like a file)
 - Socket (either a network or internal interface)
 - Message Queue/Message passing (special programming interface – passes data like internal messages/SMS)
 - Semaphore (special flag/file which controls access to resources)



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Summary

- We learnt about the difference between Programs and processes, how they are scheduled and about Inter Process Communications (IPC)
- In the linuxgym chapter 5 lab we will make use of the the IPC process called PIPES

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