

## data structures & algorithms Tutorial 10 (Week 10)







## Burning questions from the previous tutorial?

## This week's lab



### Trees

- Binary Search Tree (BST)
- Common BST methods
- BST Traversals





### Each circle is called a node



Each circle is called a node each node has two children



- Each circle is called a node each node has two children
  - - and one parent



We call all the nodes with no children leaf nodes



### Each node's children define a left subtree and a right subtree



### The top node is called the root



### The top node is called the root This is the left child of the root



The top node is called the root This is the left child of the root and this is the right child of the root











You can think of a binary tree as an extension of a linked list



You can think of a binary tree as an extension of a linked list



## struct Node { }



struct Node {
 T data {};
 Node\* left {nullptr};
 Node\* right {nullptr};
}



struct Node {
 T data {};
 Node\* left {nullptr};
 Node\* right {nullptr};
 Node\* parent {nullptr};
}



struct Node { 3 T data {}; Node\* left {nullptr}; Node\* right {nullptr}; Node Node\* parent {nullptr}; Node() {} Node(T d, Node\* node = nullptr) : data {d}, parent {node} {}



# Now time for a very important question...

## wear them



Ok now time to look at some common methods for binary search trees



### To insert an element we start at the root

### insert(7)

## 3 4 1



Now we compare if the new node is greater or less than the current node

### insert(7)

## 3 4 1



Now we compare if the new node is greater or less than the current node

### insert(7)

## 3 (1) (4)

Our new node is less than the current node, but the current node does not have a left child



### insert(7)





### So we have found the place to insert the node

### find(4)



To find a node by its key we again start at the top





### find(4)

### If the key is greater than the current node go right, otherwise go left



## Binary Search Tree

### find(4)





Bingo! If the key is equal to current node then we have found it

### find(5)





But say we were looking for 5 instead Then we would need to go right

### find(5)



### But the right child is null! So this tells us that 5 is not in the BST





min() Again we start at the root





min() But now we just always go left



## min() Once we can't go left anymore we have found the min





Starting at the root if we visit the: left child we traverse the left subtree • right child we traverse the right subtree





Preorder root  $\rightarrow$  left  $\rightarrow$  right Inorder left  $\rightarrow$  root  $\rightarrow$  right Postorder left  $\rightarrow$  right  $\rightarrow$  root





To do a preorder traversal walk along the yellow line and write down node each time you encounter a blue dot

Preorder



### Preorder [6, 3, 1, 4, 9, 8, 11]



To do an inorder traversal walk along the yellow line and write down node each time you encounter a purple dot

Inorder



### Inorder [1, 3, 4, 6, 8, 9, 11]



To do an postorder traversal walk along the yellow line and write down node each time you encounter a red dot

Postorder



### Postorder [1, 4, 3, 8, 11, 9, 6]







## Binary Search Tree

### successor(10)





## Binary Search Tree

### successor(10)

Case 1: Node has a right subtree • Find min of right subtree







[ 6 <mark>8</mark> 10 11 12 15 16 17 20 25 27 ]

## Binary Search Tree

### successor(8)





## Binary Search Tree

### successor(8)

Case 2: Node is left child • Return the parent





## Binary Search Tree

### successor(12)



![](_page_54_Figure_0.jpeg)

 $\begin{bmatrix} 6 & 8 & 10 & 11 & 12 & 15 & 16 & 17 & 20 & 25 & 27 \end{bmatrix}$ 

## Binary Search Tree

### successor(12)

Case 3: Node is right child • go to nearest ancestor for which the given node is in the left subtree

27

25