

## data structures & algorithms Tutorial 11 (Week 12)





#### Overview

- Burning questions from last week?
- This week's Lab
  - Detecting Cycles
  - DAGs
  - Topological Sort

#### This week's Lab



Today we are learning about some more useful types of graphs and graph algorithms.

#### Cycles

• Directed Acyclic Graph (DAG) • Detecting Cycles Topological Sort



#### Here we have a directed graph



#### Can anyone see any cycles in this graph?



If we start at vertex 2, can we travel along the edges and get back to vertex 2?



If we start at vertex 2, can we travel along the edges and get back to vertex 2?



If we start at vertex 2, can we travel along the edges and get back to vertex 2?



#### Yes! We have found a cycle









No! because we cannot travel from  $2 \longrightarrow 1$ . The edge is pointing the wrong way



What about if the edge is **bidirectional**? Is this a cycle?



Yes! because the edge is bidirectional so we can travel from  $2 \longrightarrow 1$ , which completes the cycle



Does this process of following edges deeper and deeper feel familiar?

What **algorithm** have we seen the past two weeks that does this?



This is very similar to DFS! So maybe we can use DFS to detect a cycle?



Let's have at look at DFS on a simple graph

Remember DFS is like exploring a maze by going down one corridor until we can't go any further and then turning back



#### We are going to keep track of the vertices that are currently part of the "corridor" we are exploring

[1]



#### [1, 2]



#### [1, 2, 3]



#### [1, 2, 3, 4]



#### [1, 2, 3]



#### [1, 2]



#### [1, 2, 5]



#### [1, 2, 5, 6]



#### [1, 2, 5]



#### [1, 2]





















#### So what would happen if we encounter a cycle along our journey?

#### [1, 2, 3]



# So what would happen if we encounter a cycle along our journey?

[1, 2, 3, 4]



#### So what would happen if we encounter a cycle along our journey?

#### [1, 2, 3]






# So what would happen if we encounter a cycle along our journey?

## [1, 2, 5]



# So what would happen if we encounter a cycle along our journey?

## [1, 2, 5, 6]



Its only during a **cycle** that we reencounter one of the vertices of a "corridor" that we are currently exploring

[1, 2, 5, 6, 2]



So we are going to use a slightly modified version of DFS to check our graph for cycles

[1, 2, 5, 6, 2]





So we are going to use a slightly modified version of DFS to check our graph for cycles

[1, 2, 5, 6, 2]



class DirectedCycle: visited = [] onStack = []

> bool containsCycle(Graph graph): bool dfs(int start, Graph graph)

So we are going to create a simple class that is going to help us find cycles in a graph

class DirectedCycle:
 visited = []
 onStack = []

bool containsCycle(Graph graph):
bool dfs(int start, Graph graph)

#### Let's start with a our basic implementation of DFS

```
void dfs(int v, Graph graph):
    visited[v] = true
    for neighbour in graph.getNeighbours(v):
        if not visited[neighbour]:
            dfs(neighbour, graph)
```

class DirectedCycle:
 visited = []
 onStack = []

bool containsCycle(Graph graph):
bool dfs(int start, Graph graph)

# We need to keep track of the current "corridor" we are exploring. So we use add the onStack vector.

### void dfs(int v, Graph graph): onStack[v] = true visited[v] = true for neighbour in graph.getNeighbours(v):

```
if not visited[neighbour]:
    dfs(neighbour, graph)
```

class DirectedCycle:
 visited = []
 onStack = []

bool containsCycle(Graph graph):
bool dfs(int start, Graph graph)

Now we also want to check if we reencounter one of the "corridor's" vertices. If we do we have found a cycle

```
bool dfs(int v, Graph graph):
    onStack[v] = true
    visited[v] = true
    for neighbour in graph.getNeighbours(v):
        if onStack[neighbour]:
            return true
        if not visited[neighbour]:
            dfs(neighbour, graph)
```

class DirectedCycle: visited = [] onStack = []

> bool containsCycle(Graph graph): bool dfs(int start, Graph graph)

```
bool dfs(int v, Graph graph):
    onStack[v] = true
    visited[v] = true
   for neighbour in graph.getNeighbours(v):
        if onStack[neighbour]:
            return true
        if not visited[neighbour]:
            existsCycle = dfs(neighbour, graph)
            if existsCycle:
                return true
```

So we now also need to update the recursive call. This call will also return true if it finds a cycle



class DirectedCycle: visited = [] onStack = []

> bool containsCycle(Graph graph): bool dfs(int start, Graph graph)

```
bool dfs(int v, Graph graph):
    onStack[v] = true
    visited[v] = true
   for neighbour in graph.getNeighbours(v):
        if onStack[neighbour]:
            return true
        if not visited[neighbour]:
            existsCycle = dfs(neighbour, graph)
            if existsCycle:
                return true
    onStack[v] = false
    return false
```

#### Finally we need to handle the case if we don't find a cycle





So far this code *kinda* works but we have some problems Can anyone see the two problems that we will run into?





































Well... the problem is there are still a lot of unexplored vertices 😕





Luckily the solution is simple, we just pick another unvisited vertex, and run the dfs algorithm on that vertex







#### So lets try dfs(5) next





#### So lets try dfs(5) next





And then we could do dfs(11) next which will visit all the vertices in that connected component





#### And then we could do dfs(9) where we will detect a cycle





Suppose that you are trying to create a learning plan for your degree so that you know what order to take your subjects in

You want to find some kind of ordering of your subjects so that you have done all the prerequisites by the time you take that class

Notice that all of the edges point in the same direction. This is called a topological sort



Also notice that we want cannot create a topological ordering if there is a cycle in our graph.

At least one edge will always point in the wrong direction

But the good news is, we just learned how to detect cycles!



Our task is to create an algorithm that can solve this problem efficiently for us!





- 1, 4, 3, 8, 11, 9, 6
  - Postorder
- Do you all remember postorder traversal?



Well here is another way to think of it. If we perform a DFS, and we write down the vertex only when we remove it from our stack

## [6] Stack

## 1, 4, 3, 8, 11, 9, 6 11) Post order

Well here is another way to think of it. If we perform a DFS, and we write down the vertex only when we remove it from our stack

## [6, 3] Stack 1, 4, 3, 8, 11, 9, 6 Post order



Well here is another way to think of it. If we perform a DFS, and we write down the vertex only when we remove it from our stack

## [6, 3, 1] Stack 1, 4, 3, 8, 11, 9, 6 Post order

4

8

Well here is another way to think of it. If we perform a DFS, and we write down the vertex only when we remove it from our stack

## [6, 3] Stack

## 1, 4, 3, 8, 11, 9, 6 11) Post order



Well here is another way to think of it. If we perform a DFS, and we write down the vertex only when we remove it from our stack

## [6, 3, 4] Stack 1, 4, 3, 8, 11, 9, 6 Post order
Well here is another way to think of it. If we perform a DFS, and we write down the vertex only when we remove it from our stack

#### [6, 3] Stack

Well here is another way to think of it. If we perform a DFS, and we write down the vertex only when we remove it from our stack

#### [6] Stack

# Topological Sort 6 $\mathbf{O}$ 8

Well here is another way to think of it. If we perform a DFS, and we write down the vertex only when we remove it from our stack

#### [ 6, 9 ] Stack

#### Topological Sort $\begin{bmatrix} 6, 9, 8 \end{bmatrix}$ 6 Stack $\mathbf{O}$ 1, 4, 3, 8, 11, 9, 6 Post order 11 8

Well here is another way to think of it. If we perform a DFS, and we write down the vertex only when we remove it from our stack

Well here is another way to think of it. If we perform a DFS, and we write down the vertex only when we remove it from our stack

### [6,9] Stack

#### Topological Sort [6, 9, 11]6 Stack $\mathbf{O}$ 1, 4, 3, 8, 11, 9, 6 Post order 11 8

Well here is another way to think of it. If we perform a DFS, and we write down the vertex only when we remove it from our stack

# Topological Sort 6 $\mathbf{O}$ 8

Well here is another way to think of it. If we perform a DFS, and we write down the vertex only when we remove it from our stack

[6, 9] Stack

Well here is another way to think of it. If we perform a DFS, and we write down the vertex only when we remove it from our stack

#### [6] Stack



Well here is another way to think of it. If we perform a DFS, and we write down the vertex only when we remove it from our stack

#### [] Stack

So let's try apply a postorder traversal to our graph of subjects



So let's try apply a postorder traversal to our graph of subjects

#### [1] Stack



So let's try apply a postorder traversal to our graph of subjects

#### [ 1, 3 ] Stack



So let's try apply a postorder traversal to our graph of subjects

[ 1, 3, 6 ] Stack



So let's try apply a postorder traversal to our graph of subjects

#### [1, 3, 6, 7] Stack



So let's try apply a postorder traversal to our graph of subjects

[1, 3, 6] Stack



So let's try apply a postorder traversal to our graph of subjects

[ 1, 3 ] Stack

7, 6 Post order



So let's try apply a postorder traversal to our graph of subjects

[1, 3, 5] Stack

> 7, 6 Post order



So let's try apply a postorder traversal to our graph of subjects

[ 1, 3 ] Stack

7, 6, 5



So let's try apply a postorder traversal to our graph of subjects

[1, 3, 4] Stack

> **7, 6, 5** Post order



So let's try apply a postorder traversal to our graph of subjects

[ 1, 3 ] Stack

7, 6, 5, 4 Post order



So let's try apply a postorder traversal to our graph of subjects

[1] Stack

7, 6, 5, 4, 3



So let's try apply a postorder traversal to our graph of subjects

[ 1, 2] Stack

7, 6, 5, 4, 3



So let's try apply a postorder traversal to our graph of subjects

[1] Stack

7, 6, 5, 4, 3, 2 Post order



So let's try apply a postorder traversal to our graph of subjects

[] Stack

7, 6, 5, 4, 3, 2, 1



So a postorder traversal gives us a reverse topological sort.

So all we need to do is reverse the result

1, 2, 3, 4, 5, 6, 7 Reverse post order



The only problem we have is we don't know what the starting vertex is.



The only problem we have is we don't know what the starting vertex is.

Thankfully this is really easy to deal with

Suppose we start at vertex 3



After performing a post order traversal starting from **vertex 3** we should have the following

#### 7, 6, 5, 4, 3 Post order



Now we just pick another unvisited vertex. Then we perform another post order traversal and we just **append** to the current result

7, 6, 5, 4, 3



Now we just pick another unvisited vertex. Then we perform another post order traversal and we just **append** to the current result

7, 6, 5, 4, 3



Now we just pick another unvisited vertex. Then we perform another post order traversal and we just **append** to the current result

7, 6, 5, 4, 3, 2 Post order



Now we just pick another unvisited vertex. Then we perform another post order traversal and we just **append** to the current result

#### 7, 6, 5, 4, 3, 2, 1 Post order



bool dfs(int v, Graph graph): onStack.insert(v) visited.insert(v) *for* neighbour in graph.getNeighbours(start): if onStack.contains(neighbour): # we have detected a cycle return true if not visited.contains(neighbour): existsCycle = dfs(neighbour, graph) *if* existsCycle: return true onStack.erase(v) order.push\_back(v) return false

```
bool computeTopoOrder(Graph graph):
 visited.clear()
for vertex in graph:
     if not visited.contains(v):
         # clear onStack for each running of dfsVisit
         onStack.clear()
         onStack.add(v)
         existsCycle = dfs(v, graph)
         if existsCycle:
             return True
 order.reverse()
 return False
```



### Student Feedback Survey

The SFS is open as of now! You can now write tutor specific feedback. Please give me some feedback. It would really mean a lot 🥰

#### https://www.sfs.uts.edu.au/

