

1. Consider the LP

$$\begin{aligned} \min z &= x_1 + x_2 \\ \text{s.t.} \quad 2x_1 - x_2 &\geq 6 \\ x_1 + 3x_2 &\leq 10 \\ x_1 - x_2 &\geq \frac{3}{2} \\ x_1, x_2 &\geq 0 \end{aligned}$$

Solve the dual problem and, with the Complementary slackness Theorem, obtain the solution for the primal problem.

2. Use dual Simplex method to solve:

$$\begin{aligned} \min z &= 5x_1 + 2x_2 + 8x_3 \\ \text{s.t.} \quad 2x_1 - 3x_2 + 2x_3 &\geq 3 \\ -x_1 + x_2 + x_3 &\geq 5 \\ x_1, x_2, x_3 &\geq 0 \end{aligned}$$

* from Linear and Non-Linear Programming by S.G.Nash and A.Sofer

Strong duality theorem

Let x be a feasible solution for primal LP (P) and y be a feasible solution of the corresponding dual LP (D). Then

$$c^T x = b^T y$$

if and only if x is an optimal solution for (P) and y is an optimal solution for (D).
(We need the dual of the LP in the standard form)

Corollary

If an optimal solution to the dual LP (D) is obtained, an optimal solution to its primal LP (P) can be readily obtained and both optimal objective values are equal.

Complementary slackness theorem

Theorem 6

- Let x be a feasible solution to the primal LP (P) and y be a feasible solution to the dual LP (D). Both solutions x and y are optimal to the primal (P) and dual (D), respectively, if and only if they satisfy

$$(c^T - y^T A)x = 0 \quad \text{and} \quad y^T (b - Ax) = 0 \quad (2)$$

THAT IS:

for each i^{th} dual
 $(c_i - y^T A_i)x_i = 0$

and

for each j^{th} primal
 $y_j (b_j - A_j x) = 0$.

constraint:

if in opt x , $x_i \neq 0$

$\downarrow$
 $c_i - y^T A_i = 0$

(active constraint)

constraint:

if in opt y , $y_j \neq 0$

$\downarrow$

$b_j - A_j x = 0$

(active constraint)

binding

1. Consider the LP

$$\begin{aligned} \min z &= x_1 + x_2 \\ \text{s.t.} \quad & 2x_1 - x_2 \geq 6 \quad \bullet \text{ normal form } \rightarrow y_1 \geq 0 \\ & x_1 + 3x_2 \leq 10 \quad \bullet \text{ reverse to normal form } \rightarrow y_2 \leq 0 \\ & x_1 - x_2 \geq \frac{3}{2} \quad \bullet \text{ normal form } \rightarrow y_3 \geq 0 \\ & x_1, x_2 \geq 0 \rightarrow \text{normal form} \rightarrow \text{dual const. in normal form} \end{aligned}$$

Solve the dual problem and, with the Complementary slackness Theorem, obtain the solution for the primal problem.

Asymmetrical dual:

$$\max w = 6y_1 + 10y_2 + \frac{3}{2}y_3$$

s.t.

$$2y_1 + y_2 + y_3 \leq 1$$

$$-y_1 + 3y_2 - y_3 \leq 1$$

$$y_1, y_3 \geq 0, \quad y_2 \leq 0$$

$$\downarrow \\ y_2' = -y_2 \rightarrow y_2' \geq 0$$

Standard form:

$$\max w = 6y_1 - 10y_2' + \frac{3}{2}y_3$$

$$\text{s.t.} \quad 2y_1 - y_2' + y_3 + s_1 = 1$$

$$-y_1 - 3y_2' - y_3 + s_2 = 1$$

$$y_1, y_2', y_3, s_1, s_2 \geq 0.$$

$$x_{B_0} = (s_1, s_2)$$

max problem!

	y_1	y_2'	y_3	s_1	s_2	RHS
w	-6	10	$-\frac{3}{2}$	0	0	0
s_1	2	-1	1	1	0	1
s_2	-1	0-3	-1	0	1	1
w	0	7	$\frac{3}{2}$	3	0	3
y_1	1	$-\frac{1}{2}$	$\frac{1}{2}$	$\frac{1}{2}$	0	$\frac{1}{2}$
s_2	0	$-\frac{7}{2}$	$-\frac{1}{2}$	$\frac{1}{2}$	1	$\frac{3}{2}$

$$R_0' = R_0 + 6R_1'$$

$$R_1' = R_1 / 2$$

$$R_2' = R_2 + R_1'$$

optimal $\hat{C}_n > 0$

$$\left. \begin{aligned} w^* &= 3 \\ y &= \begin{pmatrix} 1/2 \\ 0 \\ 0 \end{pmatrix} \end{aligned} \right\} \begin{array}{l} \text{dual opt. solution and} \\ \text{OF value} \end{array}$$

$$\downarrow \\ z^* = w^* = 3 \quad \text{by Strong Duality Theorem}$$

$$\downarrow \\ x_1^* + x_2^* = 3$$

as $y_1 \neq 0 \rightarrow 1^{st} (P)$ constraint is active "="
by comp. slackness Theorem

$$\downarrow \\ 2x_1^* - x_2^* = 6$$

$$\downarrow \\ \begin{cases} x_1^* + x_2^* = 3 \\ 2x_1^* - x_2^* = 6 \end{cases} + \\ \hline 3x_1^* = 9 \rightarrow x_1^* = 3 \\ x_2^* = 0$$

2. Use dual Simplex method to solve:

$$\min z = 5x_1 + 2x_2 + 8x_3$$

$$\text{s.t. } 2x_1 - 3x_2 + 2x_3 \geq 3$$

$$-x_1 + x_2 + x_3 \geq 5$$

$$x_1, x_2, x_3 \geq 0$$

(P) $\rightarrow$ (D)

$$\max 3y_1 + 5y_2$$

$$\text{s.t. } 2y_1 - y_2 \leq 5$$

$$-3y_1 + y_2 \leq 2$$

$$2y_1 + y_2 \leq 8$$

$$y_1, y_2 \geq 0$$

* from Linear and Non-Linear Programming by S.G.Nash and A.Soffer

Standard form ()

$$\min z = 5x_1 + 2x_2 + 8x_3$$

$$\text{s.t. } 2x_1 - 3x_2 + 2x_3 - e_1 = 3$$

$$-x_1 + x_2 + x_3 - e_2 = 5$$

$$x_1, x_2, x_3, e_1, e_2 \geq 0.$$

not canonical
+ (-1)

$$x_B = (e_1, e_2)$$

$$-2x_1 + 3x_2 - 2x_3 + e_1 = -3$$

$$x_1 - x_2 - x_3 + e_2 = -5$$

$$x_1, x_2, x_3, e_1, e_2 \geq 0.$$

$$\min z = 5x_1 + 2x_2 + 8x_3$$

$$z - 5x_1 - 2x_2 - 8x_3 = 0$$

$$\text{or } \hat{C}_N^T = C_B^T B_0^{-1} N - C^T =$$

$$z(x_{B_0}) = C_B^T B_0^{-1} b = 0$$

$$-C^T = (-5, -2, -8)$$

	x_1	x_2	x_3	e_1	e_2	RHS
z	-5	-2	-8	0	0	0
e_1	-2	3	-2	1	0	-3 < 0
e_2	1	-1	-1	0	1	-5 < 0
z	-7	0	-6	0	-2	10
e_1	1	0	-5	1	3	-18 < 0
x_2	-1	1	1	0	-1	5
z	$-\frac{41}{5}$	0	0	$-\frac{6}{5}$	$-\frac{28}{5}$	$\frac{158}{5}$
x_3	$-\frac{1}{5}$	0	1	$-\frac{1}{5}$	$-\frac{3}{5}$	$\frac{18}{5}$
x_2	$-\frac{4}{5}$	1	0	$\frac{1}{5}$	$-\frac{2}{5}$	$\frac{7}{5}$

① Optimal, not Feasible

Ratio test:

$$\left| \frac{-2}{-1} \right| < \left| \frac{-8}{-1} \right|$$

e_2 leaving
as $-5 < -3$ x_2 enters

$$R'_0 = R_0 + 2R'_2$$

$$R'_1 = R_1 - 3R'_2$$

② Optimal, not Feasible

$$R'_2 = R_2 \times (-1)$$

$$R'_0 = R_0 + 6R'_1$$

$$R'_1 = \frac{R_1}{(-5)}$$

$$R'_2 = R_2 - R'_1$$

Optimal as $\hat{C}_N < 0$

$$z^* = \frac{158}{5} = 31.6$$

$$x = \begin{pmatrix} 0 \\ 7/5 \\ 18/5 \end{pmatrix} \neq 0 \rightarrow 2^{nd}, 3^{rd} \text{ dual const. are active}$$

Dual

$$\begin{aligned} \max \quad & 3y_1 + 5y_2 \\ \text{s.t.} \quad & 2y_1 - y_2 \leq 5 \\ & -3y_1 + y_2 \leq 2 \\ & 2y_1 + y_2 \leq 8 \\ & y_1, y_2 \geq 0 \end{aligned}$$

$$w^* = z^* = \frac{158}{5} \rightarrow \begin{cases} 3y_1^* + 5y_2^* = \frac{158}{5} \\ -3y_1^* + y_2^* = 2 \end{cases} \rightarrow \begin{aligned} 6y_2^* &= \frac{168}{5} \rightarrow y_2^* = \frac{168}{6 \times 5} = \frac{28}{5} \\ y_1^* &= \frac{6}{5} \end{aligned}$$